

Example x86 Code

Computer Systems Sections 3.3 - 3.6

Simple C Function

```
int myfunc(int a,int b) {  
    int c;  
    c = a + 3;  
    c = c + b;  
    return c;  
}
```

gcc -m32 -O0 -S prog.c

prog.c

```
int myfunc(int a,int b) {  
    int c;  
    c = a + 3;  
    c = c + b;  
    return c;  
}
```

prog.s

myfunc:

```
pushl    %ebp  
movl     %esp, %ebp  
subl     $16, %esp  
movl     8(%ebp), %eax  
addl     $3, %eax  
movl     %eax, -4(%ebp)  
movl     12(%ebp), %eax  
addl     %eax, -4(%ebp)  
movl     -4(%ebp), %eax  
leave  
ret
```

Function Invocation Details (Lecture 10)

[OS(caller) invokes Main (callee)]
[Main (callee) startup]

a=3;
b=4;

- Evaluate arguments
- Copy argument values to parameters
- Preserve Caller's state
- Save return address

[Main (caller) invokes myfn(callee)]
 x=a, y=b+2;

[myfn (callee) startup]

- Preserve caller's state
- Create/initialize local variables
- Establish argument accessibility

x*x + y*y

[myfn (callee) return to main]

[main (caller) return from myfn]

c = [return value from myfn]

[main (callee) return to OS]

myfunc entrance

- Preserve caller's state
- Create/initialize local variables
- Establish argument accessibility

```
int myfunc(int a,int b) {
    int c;
    c = a + 3;
    c = c + b;
    return c;
}
```

myfunc:

```
pushl %ebp
movl %esp, %ebp
subl $16, %esp
movl 8(%ebp), %eax
addl $3, %eax
movl %eax, -4(%ebp)
movl 12(%ebp), %eax
addl %eax, -4(%ebp)
movl -4(%ebp), %eax
leave
ret
```

Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	xdead beef
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 0000

myfunc statement 1

myfunc:

```
int myfunc(int a,int b) {
    int c;
    c = a + 3;
    c = c + b;
    return c;
}
```

```
pushl %ebp
movl %esp, %ebp
subl $16, %esp
movl 8(%ebp), %eax
addl $3, %eax
movl %eax, -4(%ebp)
movl 12(%ebp), %eax
addl %eax, -4(%ebp)
movl -4(%ebp), %eax
leave
ret
```

Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	xdead beef
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 0003

myfunc statement 1

myfunc:

```
int myfunc(int a,int b) {
    int c;
    c = a + 3;
    c = c + b;
    return c;
}
```

```
pushl %ebp
movl %esp, %ebp
subl $16, %esp
movl 8(%ebp), %eax
addl $3, %eax
movl %eax, -4(%ebp)
movl 12(%ebp), %eax
addl %eax, -4(%ebp)
movl -4(%ebp), %eax
leave
ret
```

Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	xdead beef
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 0006

myfunc statement 1

myfunc:

```
int myfunc(int a,int b) {
    int c;
    c = a + 3;
    c = c + b;
    return c;
}
```

```
pushl   %ebp
movl    %esp, %ebp
subl    $16, %esp
movl    8(%ebp), %eax
addl    $3, %eax
movl    %eax, -4(%ebp)
movl    12(%ebp), %eax
addl    %eax, -4(%ebp)
movl    -4(%ebp), %eax
leave
ret
```

Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	x0000 0006
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 0006

myfunc statement 2

myfunc:

```
int myfunc(int a,int b) {
    int c;
    c = a + 3;
    c = c + b;
    return c;
}
```

```
pushl %ebp
movl %esp, %ebp
subl $16, %esp
movl 8(%ebp), %eax
addl $3, %eax
movl %eax, -4(%ebp)
movl 12(%ebp), %eax
addl %eax, -4(%ebp)
movl -4(%ebp), %eax
leave
ret
```

Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	x0000 0006
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 0004

myfunc statement 2

myfunc:

```
int myfunc(int a,int b) {
    int c;
    c = a + 3;
    c = c + b;
    return c;
}
```

```
pushl %ebp
movl  %esp, %ebp
subl  $16, %esp
movl  8(%ebp), %eax
addl  $3, %eax
movl  %eax, -4(%ebp)
movl  12(%ebp), %eax
addl  %eax, -4(%ebp)
movl  -4(%ebp), %eax
leave
ret
```

Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	x0000 000A
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 0004

Function Return Details

[OS(caller) invokes Main (callee)]
[Main (callee) startup]

a=3;
b=4;

[Main (caller) invokes myfn(caller)]
x=a, y=b+2;

[myfn (callee) startup]

x*x + y*y

[myfn (callee) return to main]

[main (caller) return from myfn]

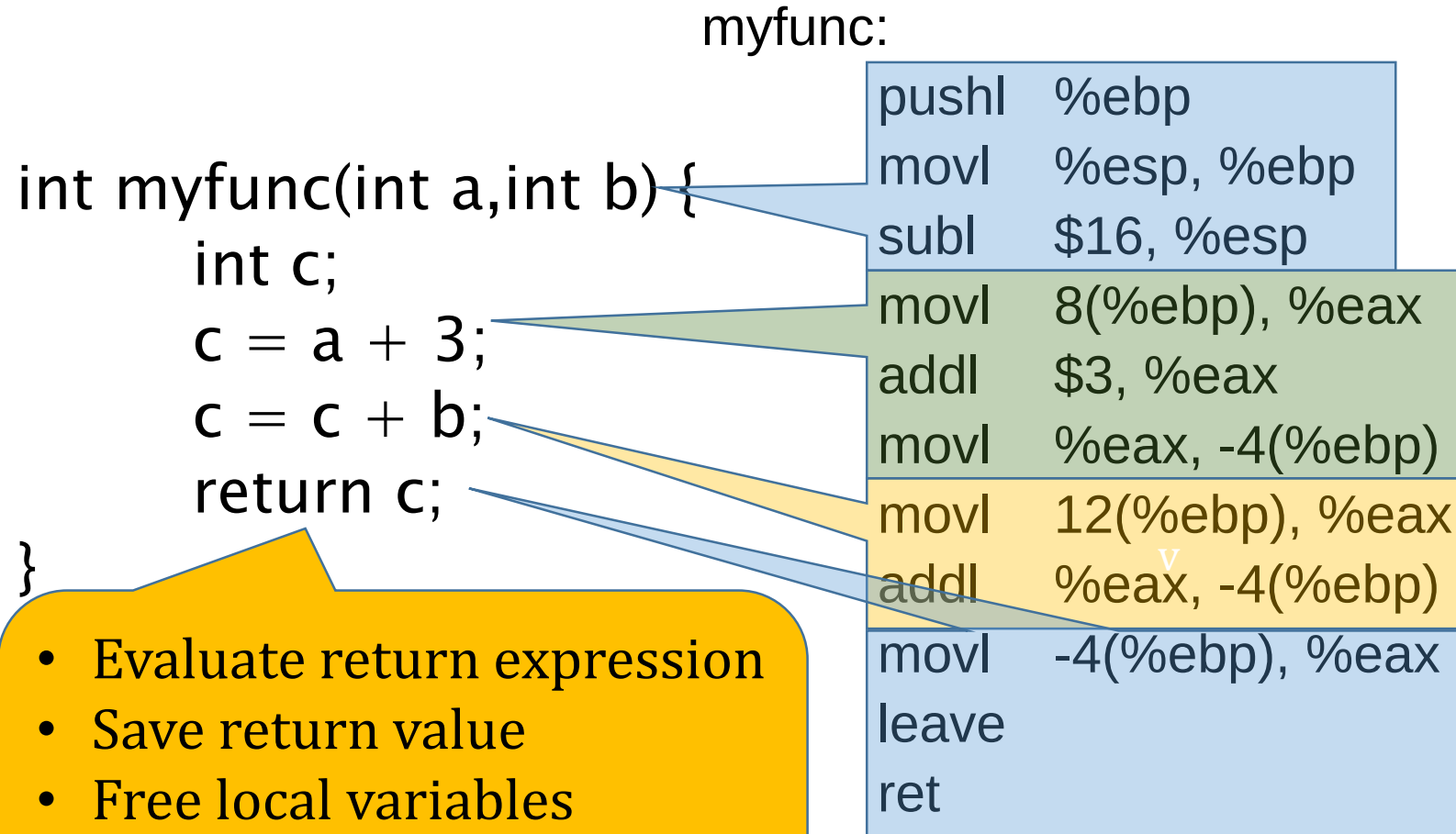
c = [return value from myfn]

[main (callee) return to OS]

- Evaluate return expression
- Save return value
- Free local variables
- Restore caller's state
- Branch to return address

- Restore caller's state
- Free space for arguments/return address
- Use returned value

myfunc return



Memory		
	xFFFF FFF0	
b	xFFFF FFEC	x0000 0004
a	xFFFF FFE8	x0000 0003
	xFFFF FFE4	
	xFFFF FFE0	
c	xFFFF FFDC	x0000 000A
	xFFFF FFD8	
	...	

Reg	Value
ebp	xFFFF FFE0
eax	x0000 000A

- Evaluate return expression
- Save return value
- Free local variables
- Restore caller's state
- Branch to return address

If/Then/Else

```
int myfunc(int a,int b) {
    int c;
    if (a>b) c = a;
    else c = b;
    return c;
}
```

```

pushl   %ebp
movl    %esp, %ebp
subl    $16, %esp
movl    8(%ebp), %eax
cmpl    12(%ebp), %eax
jle     .L2
movl    8(%ebp), %eax
movl    %eax, -4(%ebp)
jmp     .L3

.L2:
movl    12(%ebp), %eax
movl    %eax, -4(%ebp)

.L3:
movl    -4(%ebp), %eax
leave
ret
```

```
int myfunc(int a,int b) {  
    int c;  
    if (a>b)  
        c = a;  
    else  
        c = b;  
    return c;  
}
```

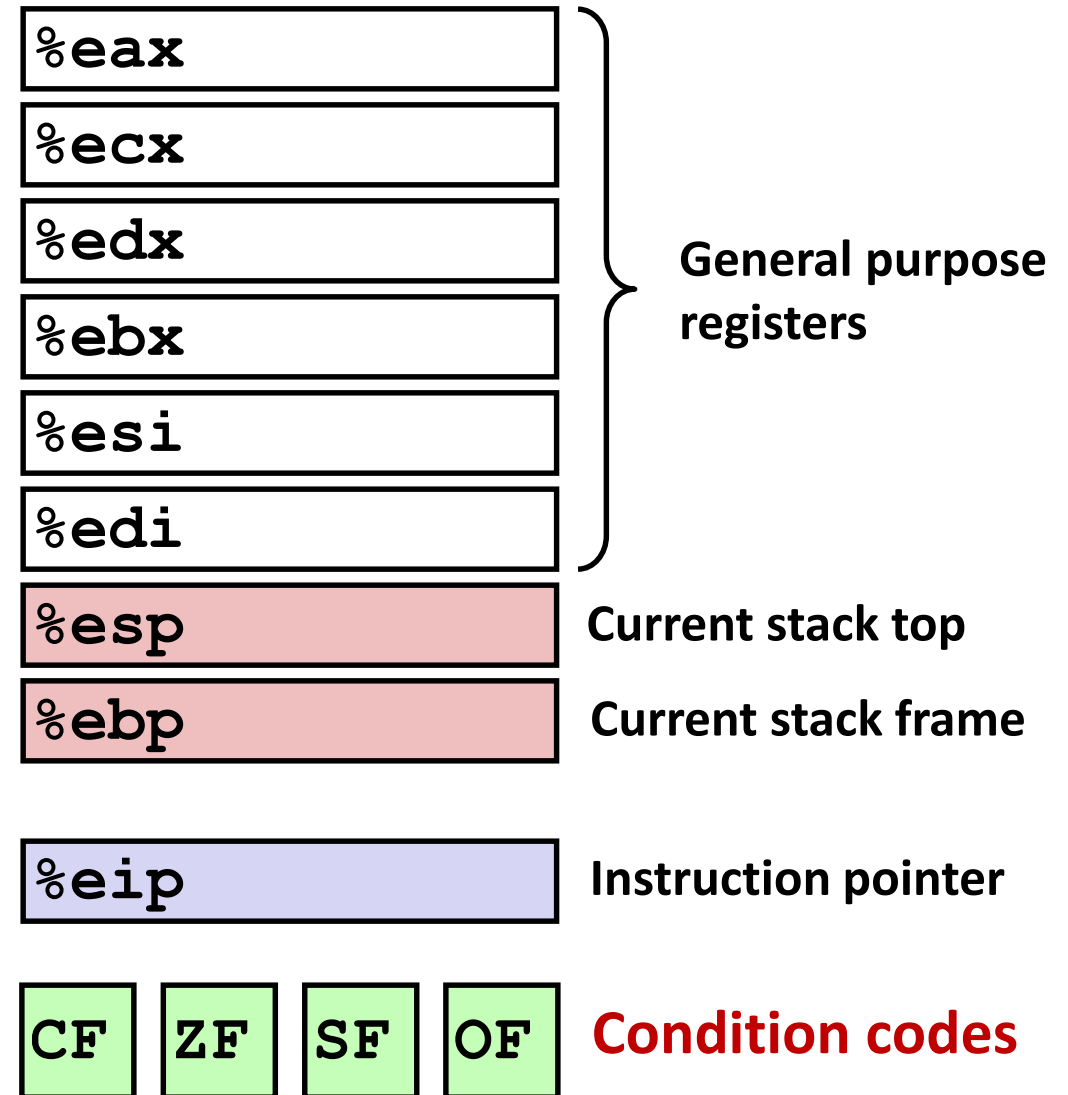
```
movl    -4(%ebp), %eax
leave
ret
```

[illegible]

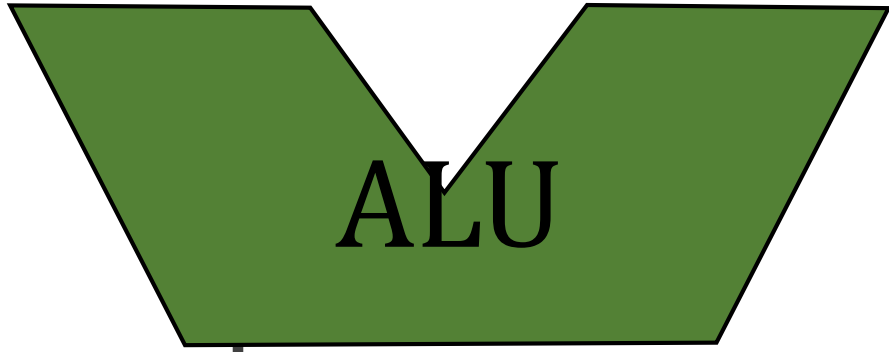
Processor State (IA32)

Information about currently
executing program

- Temporary data
(**%eax**, ...)
- Location of runtime stack
(**%ebp**, **%esp**)
- Location of current code control
point
(**%eip**, ...)
- Status of recent tests
(**CF**, **ZF**, **SF**, **OF**)
- ...



Condition Code Registers



- **CF** Carry Flag = 1 if last result MSB overflow (unsigned)
- **ZF** Zero Flag = 1 if last result was zero
- **SF** Sign Flag = 1 if last result < 0 (MSB==1)
- **OF** Overflow Flag = 1 if last result sign is incorrect (++- or --+)

Condition Codes (Implicit Setting)

- Implicitly set by arithmetic operations

Example: `addl a,b ; b' = a + b`

CF set if carry out from most significant bit (unsigned overflow)

ZF set if `b' == 0`

SF set if `b' < 0` (as signed)

OF set if two's-complement (signed) overflow

`(a > 0 && b > 0 && b' < 0) || (a < 0 && b < 0 && b' >= 0)`

- Not set by `leal` instruction
- [Full documentation](#) (IA32) or [Wikibooks X86 Control Flow](#)

Condition Codes (Explicit Set: Compare)

Condition code can be set explicitly by the “cmp” instruction...

- **cmp_l / cmp_q** *Src2, Src1*
- **cmp_l** **b**, **a** like computing **a-b** without setting destination
- **ZF set** if **a == b**
- **SF set** if **(a-b) < 0** (as signed)
- **OF set** if two's-complement (signed) overflow
(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)
- **CF set** if carry out from most significant bit (used for unsigned comparisons)

Condition Code Interpretations

Based on $\text{CMP } B, A \rightarrow A - B$ or $\text{SUB } B, A \rightarrow A = A - B$

Signed Comparisons

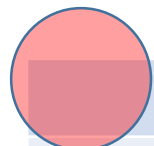
Branch	Means	ZF	SF	OF	CF
je	$a == b$	1	X	X	X
jne	$a != b$	0	X	X	X
jg	$a > b$	0	0	0	X
		0	1	1	X
jge	$a \geq b$	X	0	0	X
		X	1	1	X
jl	$a < b$	0	1	0	X
		0	0	1	X
jle	$a \leq b$	X	1	0	X
		X	0	1	X

Unsigned Comparisons

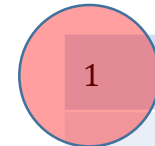
Branch	Means	ZF	SF	OF	CF
je	$a == b$	1	X	X	X
jne	$a != b$	0	X	X	X
ja	$a > b$	0	X	X	0
jae	$a \geq b$	X	X	X	0
jb	$b > a$	X	X	X	1
jbe	$b \geq a$	X	X	X	1
		1	X	X	X

Carry Flag Confusion

- Question: How does the carry flag get set when performing unsigned subtraction?
- Official Answer: The carry flag is set if either the most significant bit overflows OR if a borrow is required to the most significant bit to perform the subtraction



							1
	0	...	0	0	0	1 0	0
-	0	...	0	0	0	0	1
	1	...	1	1	1	0	1



1	1	...	1	1	1	1	
	1 0	...	1 0	1 0	1 0	0	1
-	0	...	0	0	0	1	0
	1	...	1	1	1	0	1

Carry Flag Confusion

- Unofficial Answer: The carry flag is set if either the most significant bit overflows OR if $A + (-B)$ results in negative (assuming extra sign bit)

1	1	1	...	1	1	1		
	0	0	...	0	0	0	1	0
+	1	1	...	1	1	1	1	1
	0	0	...	0	0	0	0	1

	SGN							
	0	0	...	0	0	0	0	1
+	1	1	...	1	1	1	1	0
	1	1	...	1	1	1	0	1

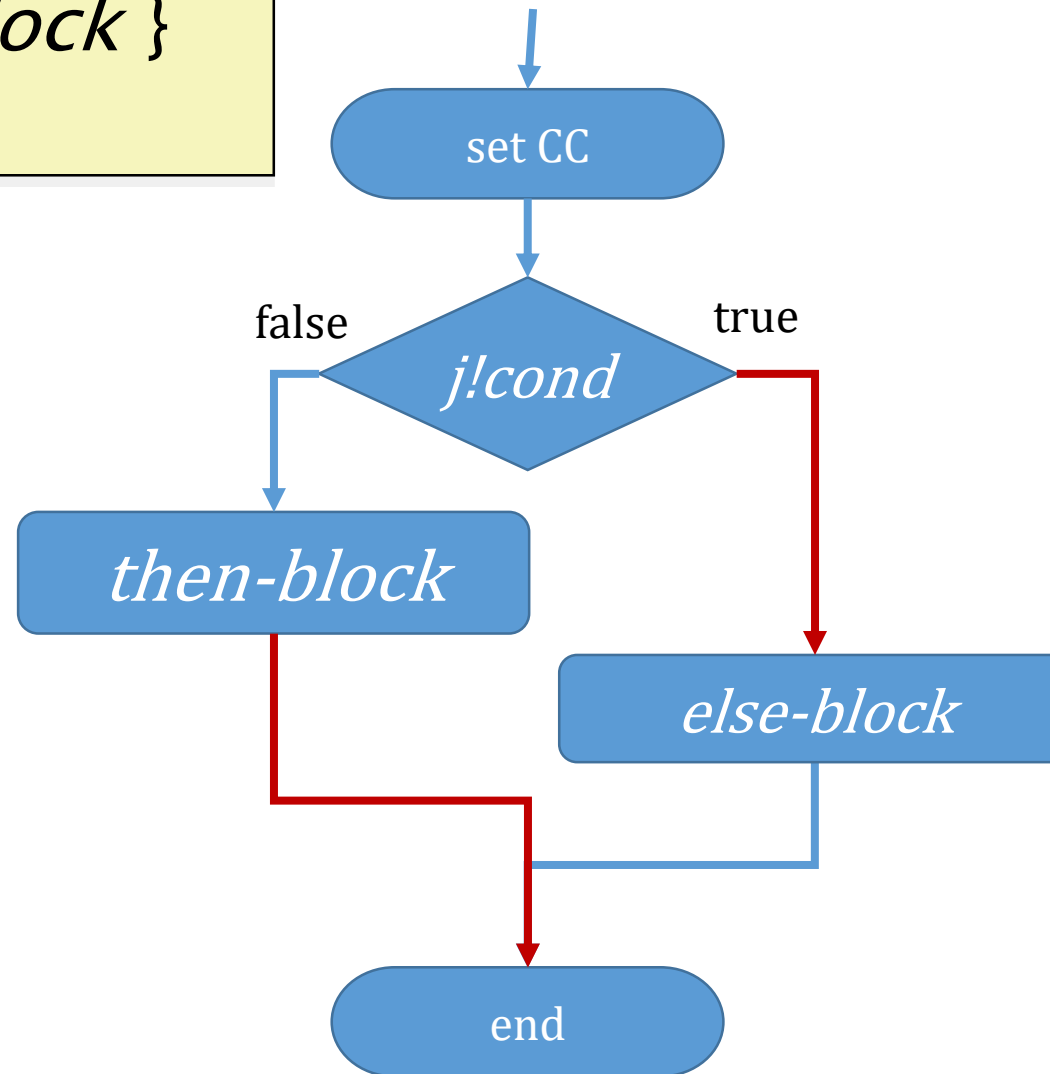
Condition Codes (Explicit Set: Test)

Condition code can be set explicitly by the “test” instruction...

- `test Src2, Src1`
- `testl a, b` like computing `a&b` without setting destination
- **ZF set** if `a & b == 0` (a bitwise different from b)
- **SF set** if `(a&b)` most significant bit=1 (sign bit)
- **OF unset (0)**
- **CF unset (0)**

Translating C to x86: If/then/else

```
if ( cond ) { then-block }
else { else-block }
```



```

...
cmp1    ...

jxx     .L6

...      ; then block
jmp     .L7

.L6:
...      ; else block

.L7:
```

```
int myfunc(int a,int b) {  
    int c;  
    if (a>b) c = a;  
    else c = b;  
    return c;  
}
```

```
movl    -4(%ebp), %eax
leave
ret
```

[illegible]